

Data is Flowing in the Wind: A Review of Data-Flow Integrity Methods to Overcome Non-Control-Data Attacks

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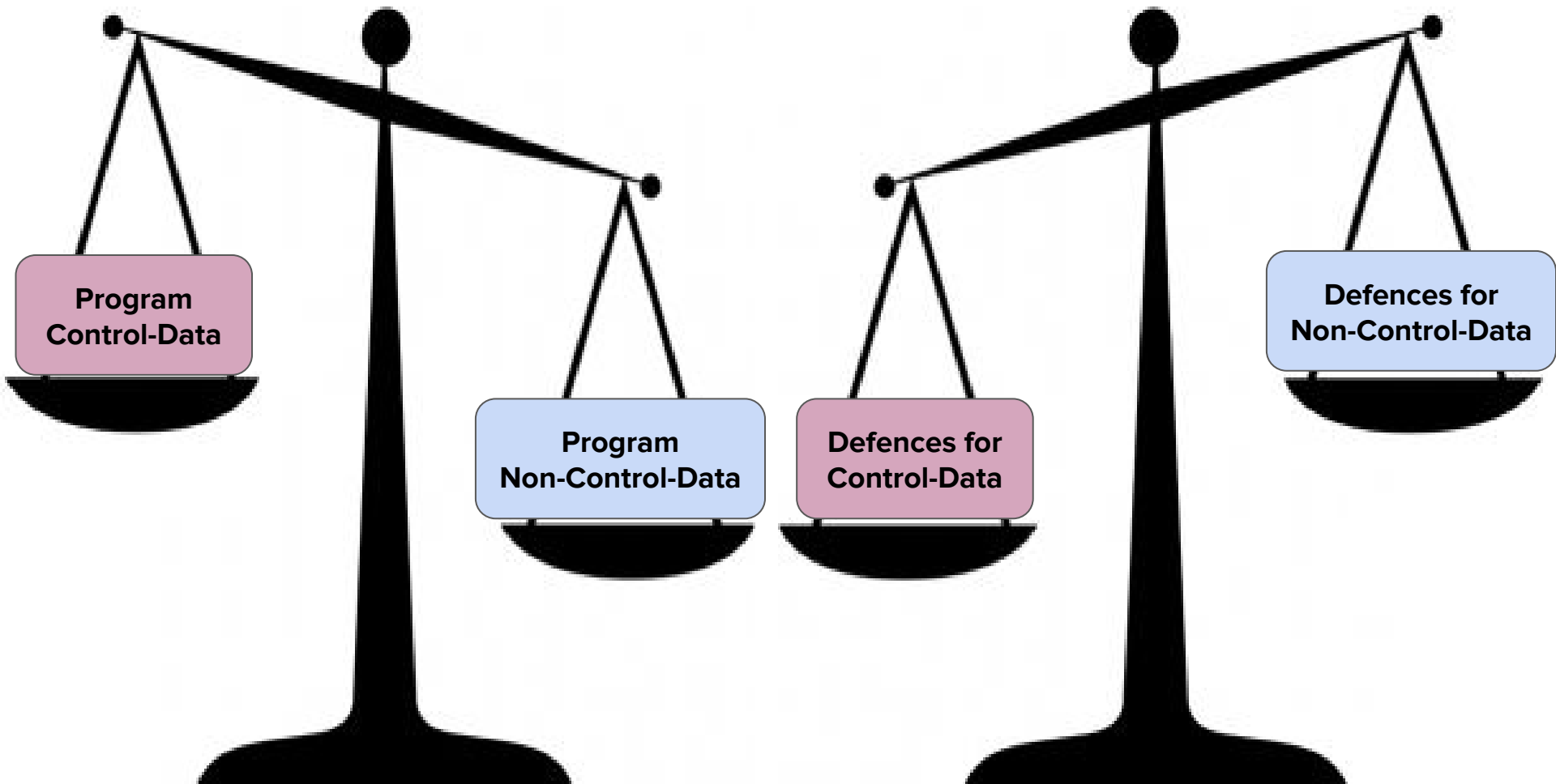
Rationale



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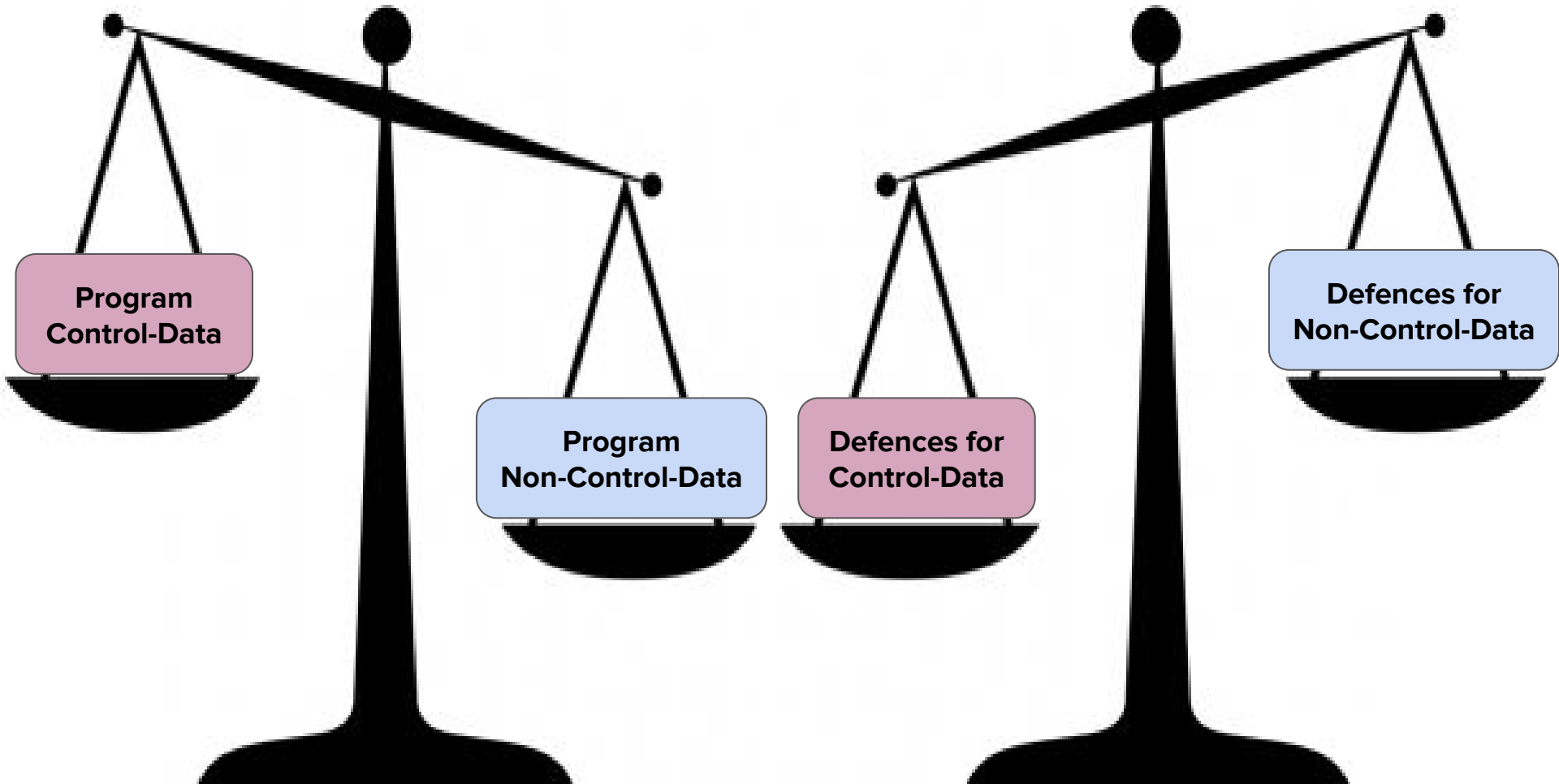


Rationale



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Maybe I should switch to attack non-control data?



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Maybe I should switch to attack non-control data?



I can sense a storm coming....



**Program
Control-Data**

**Program
Non-Control-Data**

**Defences for
Control-Data**

**Defences for
Non-Control-Data**

Rationale

Maybe I should switch to attack non-control data?



Non-Control-Data Attacks Are Realistic Threats.
Chen et al. Usenix Sec'05

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Program Control-Data

Program Non-Control-Data

Defences for Control-Data

Defences for Non-Control-Data

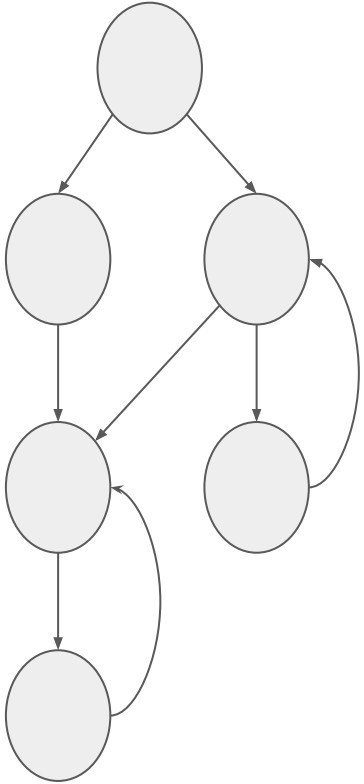
Control Data VS Non-Control-Data Attacks

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Control-Data Attacks

Modify the control-flow of a program

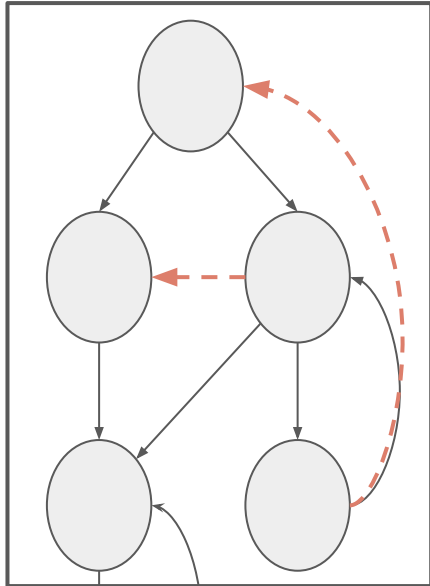
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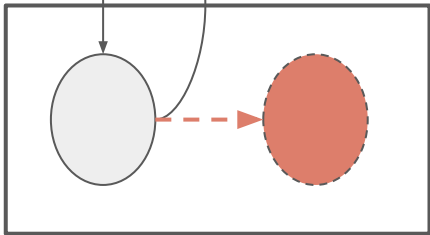
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Control Data VS Non-Control-Data Attacks



Code reuse

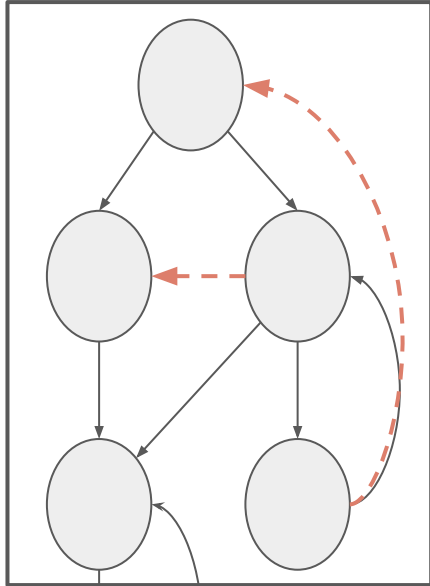


Code injection

Control-Data Attacks

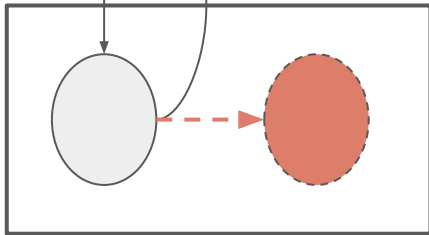
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Control Data VS Non-Control-Data Attacks



Code reuse

Modify the values of
`ret`, `call`, `jmp`
instructions

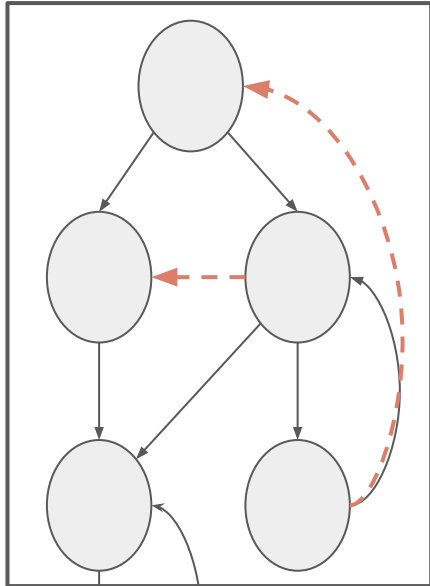


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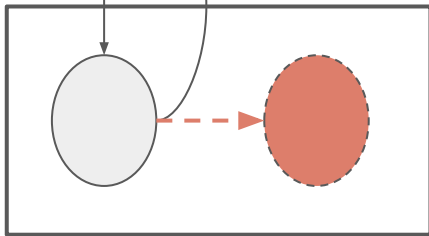
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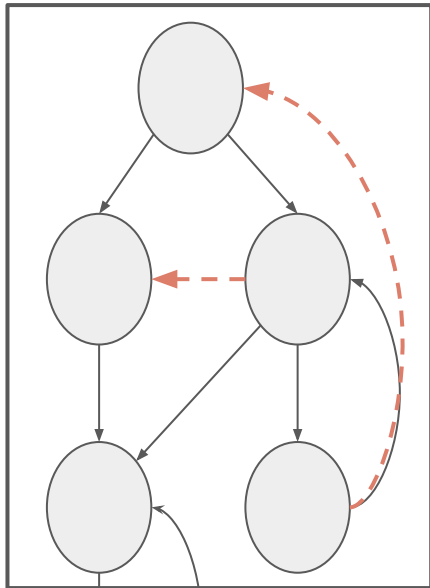
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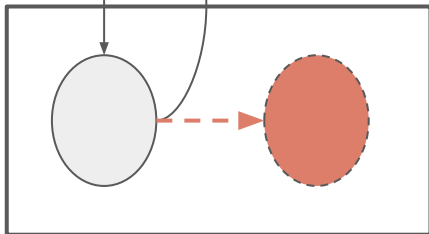
Do not affect the control-flow of a
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Control Data VS Non-Control-Data Attacks

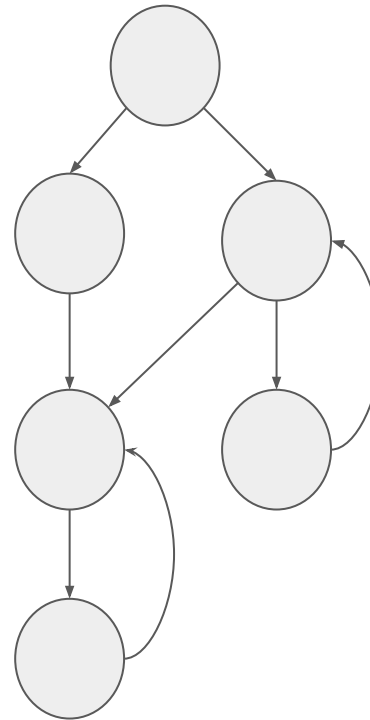


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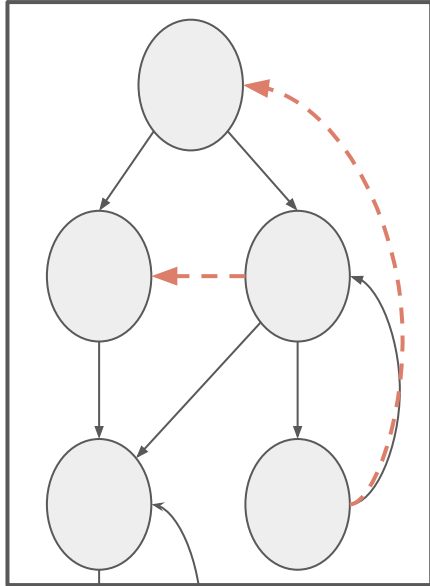
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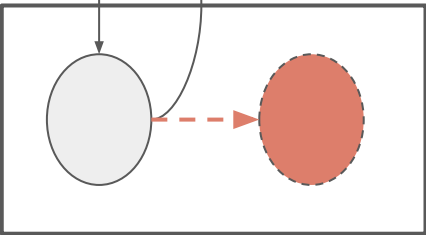
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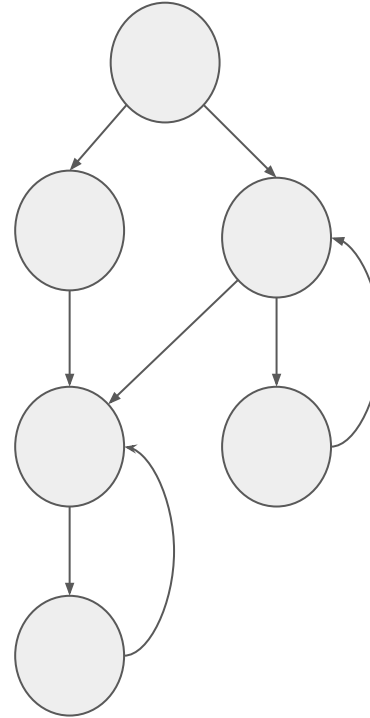


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Code injection



Remain invisible to
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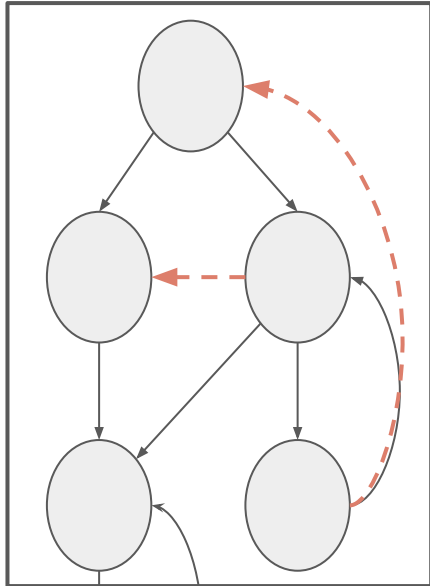
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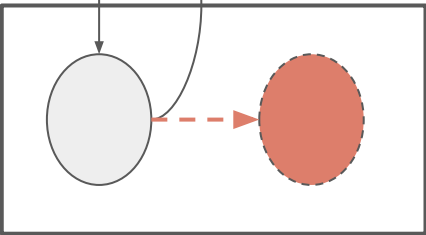
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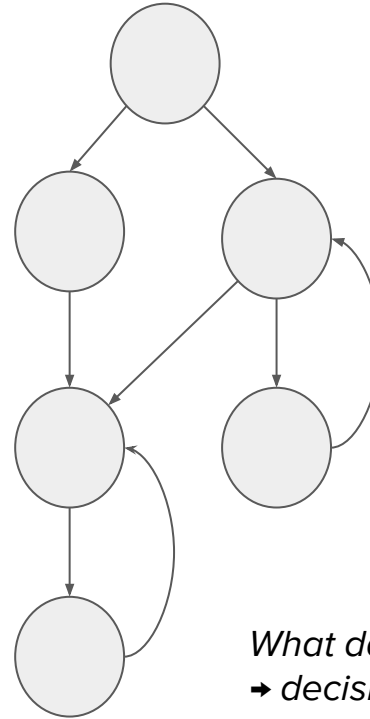


Code reuse

Modify the values of
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Code injection



Remain invisible to
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What data do they target?
→ *decision making data,*
user input etc.

Control-Data Attacks

Modify the control-flow of a
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Non-Control-Data Attacks

Do not affect the control-flow of a
program

Security-Critical Non-Control Data

Chen et al. Usenix Sec'05

Hu et al. Usenix Sec'15

Security-Critical Non-Control Data

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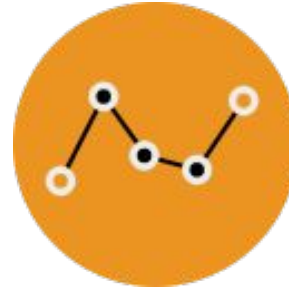
Hu et al. Usenix Sec'15



**Configuration
data**



**User identity
data**



**Decision-making
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**User input
data**

Security-Critical Non-Control Data

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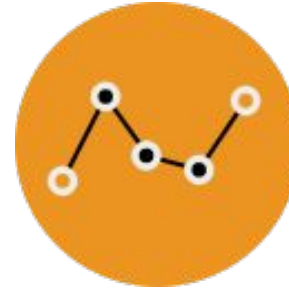
Hu et al. Usenix Sec'15



**Configuration
data**



**User identity
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**Decision-making
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**User input
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**Passwords & private
keys**



**System call
parameters**



Randomised values

A sample non-control-data attack

A sample non-control-data attack

```
struct passwd
{
    uid_t pw_uid;
    ...
} *pw;
...
int uid = getuid();
pw->pw_uid = uid; // save current uid
// [ format string vulnerability ]
...
void passive(void)
{
    ...
    setuid(0); // become root
    ...
setuid(pw->pw_uid); // drop root priv
}
```

From wu-ftpd web server.
Hu et al. Usenix Sec'15

A sample non-control-data attack

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struct passwd
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void passive(void)
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}
```

```
int uid = getuid();
pw->pq_uid = uid; // save current uid
// [ format string vulnerability ]
// [ exploit it to overwrite 'uid' ]
// [ pw->pq_uid = 0; ]
void passive(void)
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    seteuid(pw->pw_uid); // avoid priv. drop
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Hu et al. Usenix Sec'15

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Circumvents Control
Flow Integrity?



Data-Flow Stitching

Data-Flow Stitching

Hu et al. Usenix Sec'15

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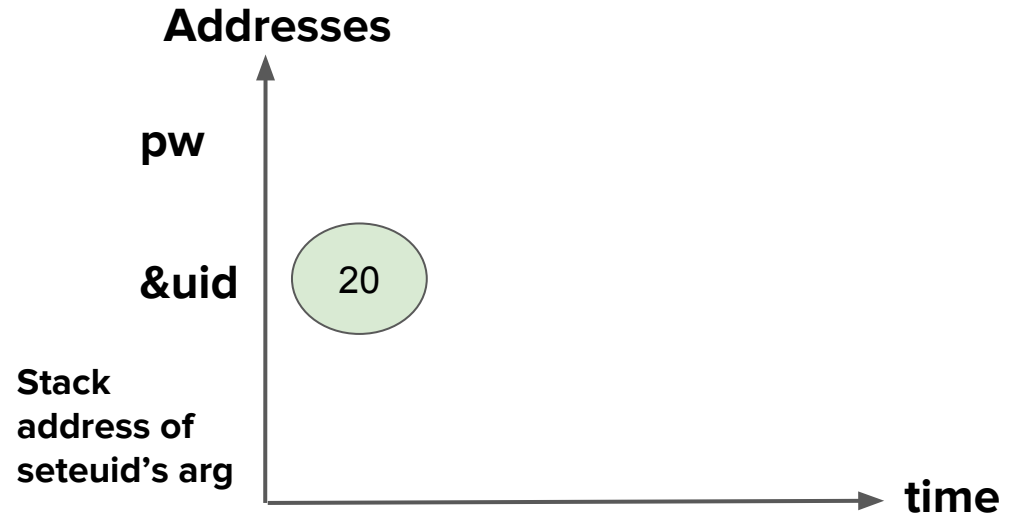
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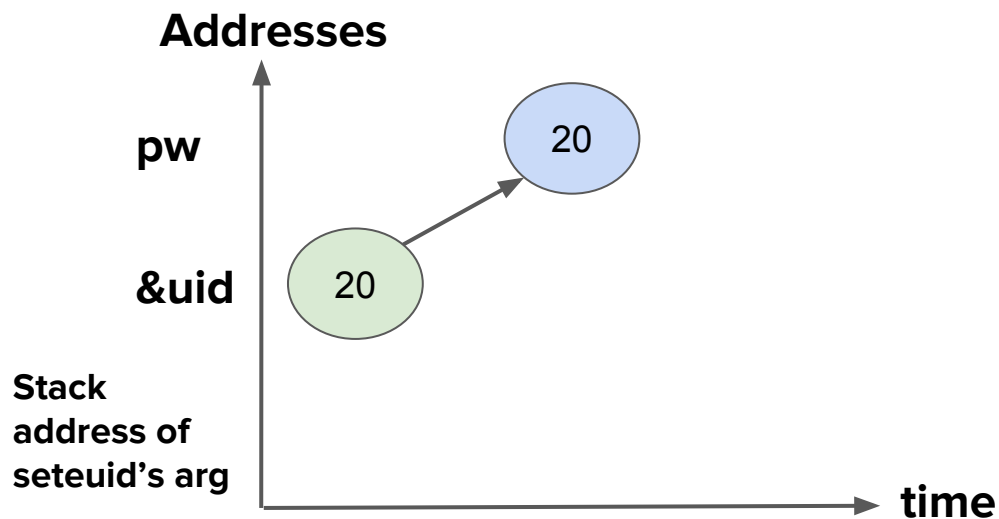
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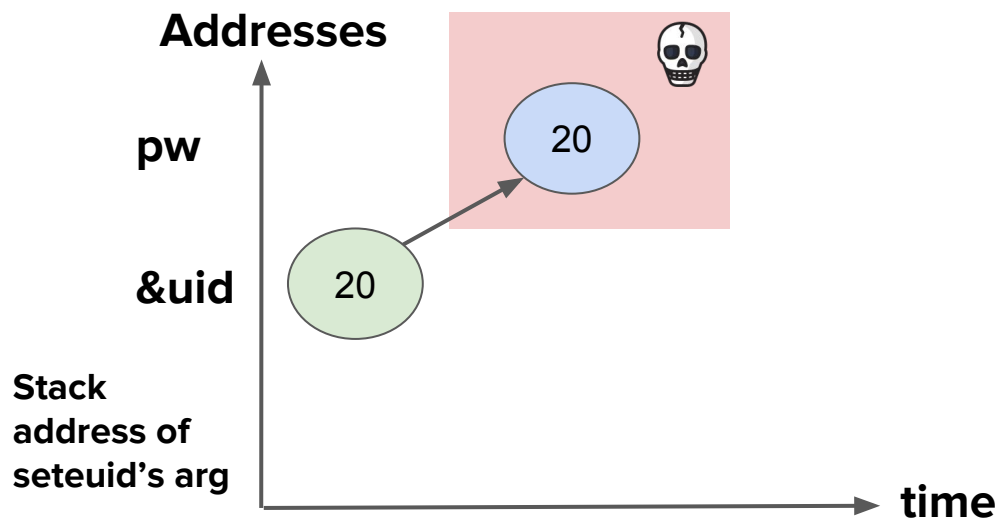
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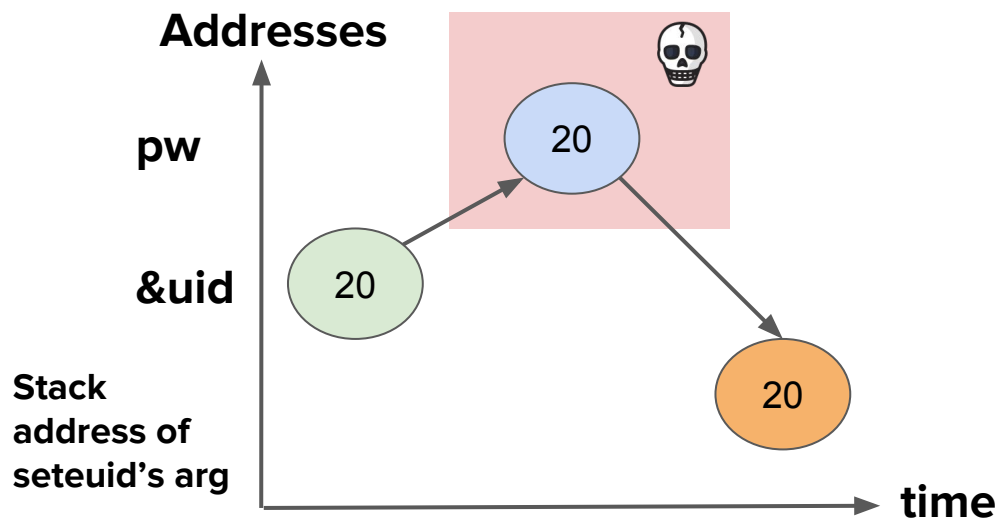
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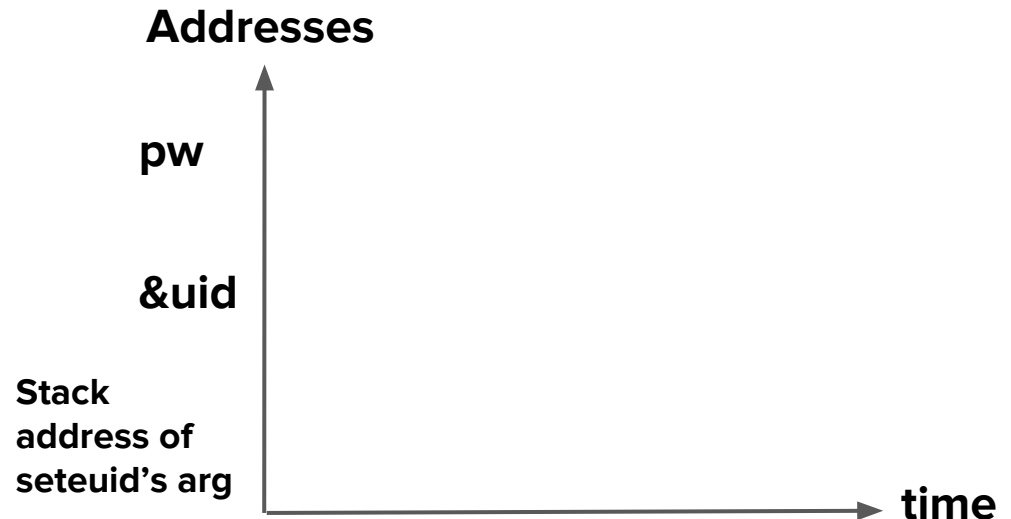
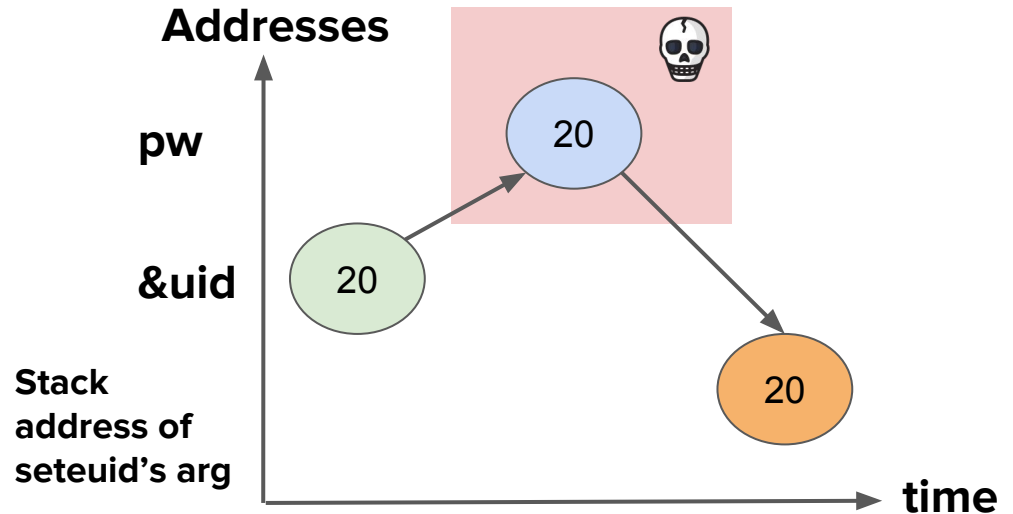
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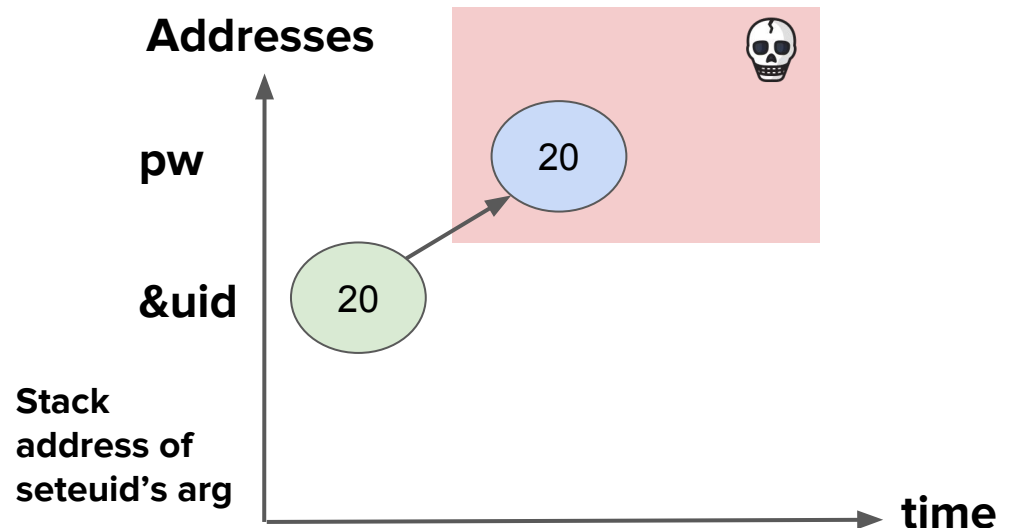
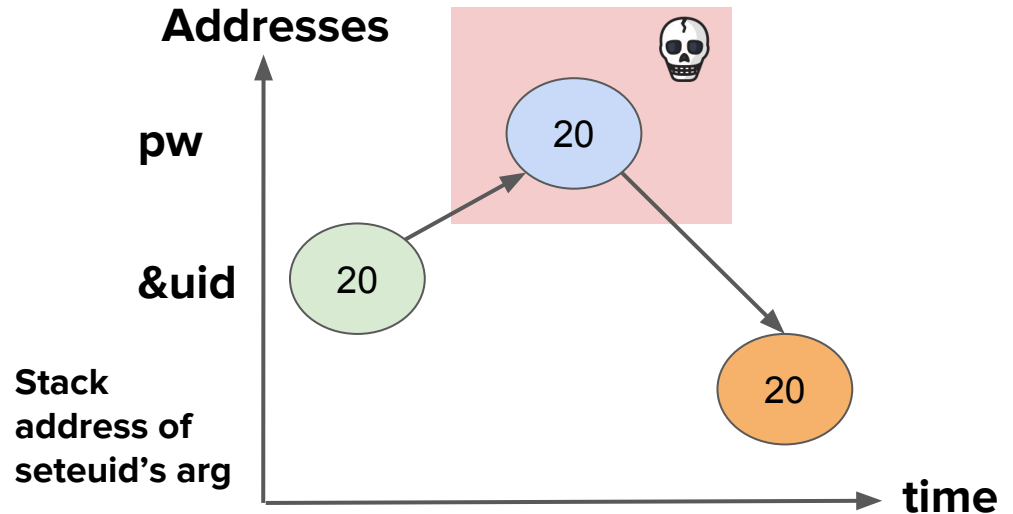
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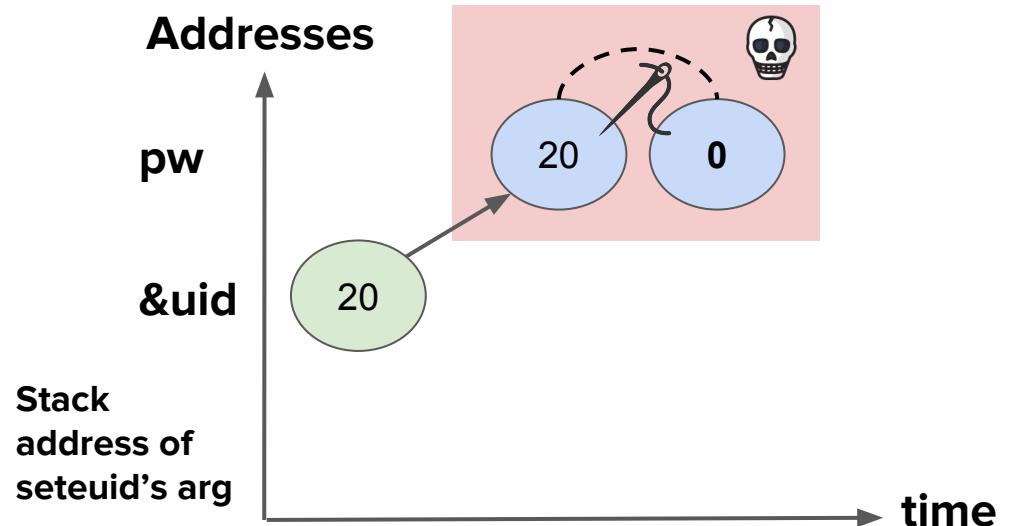
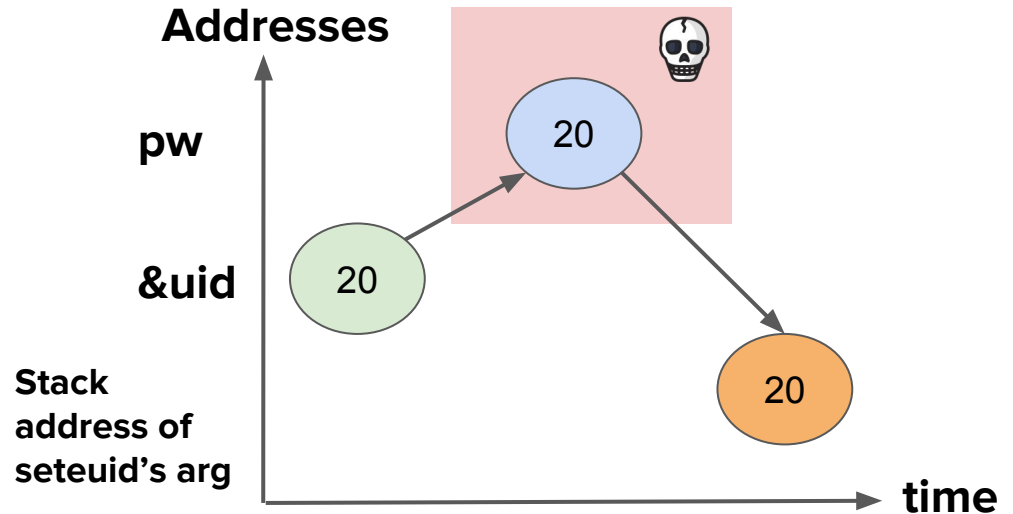
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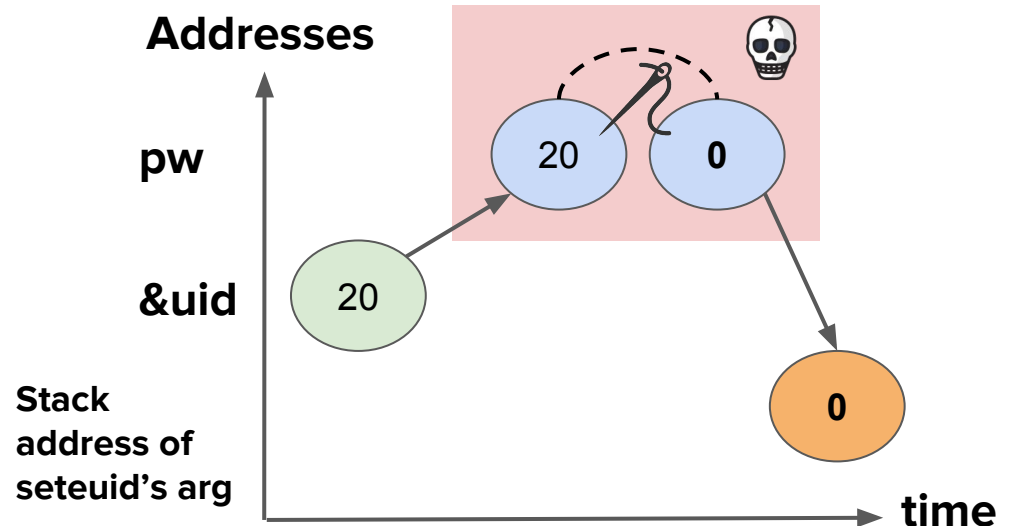
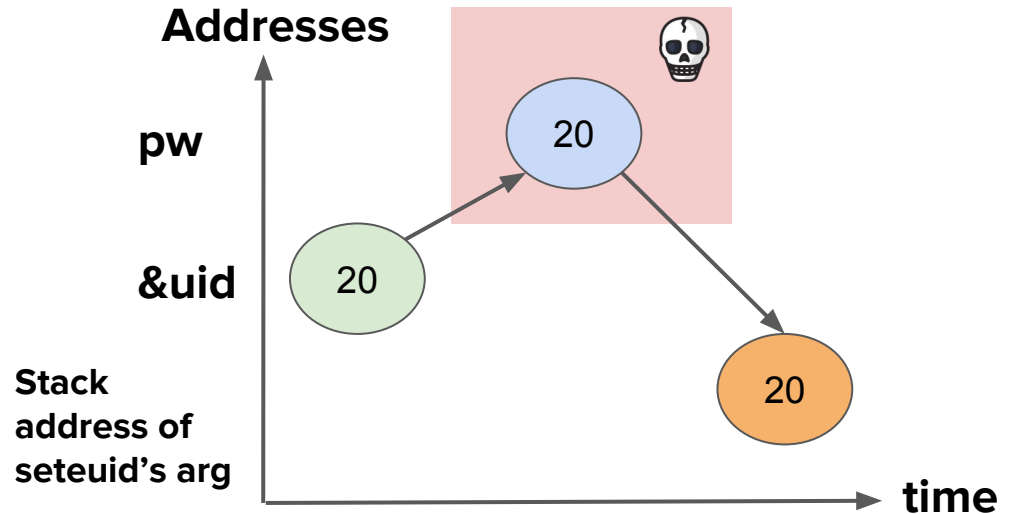
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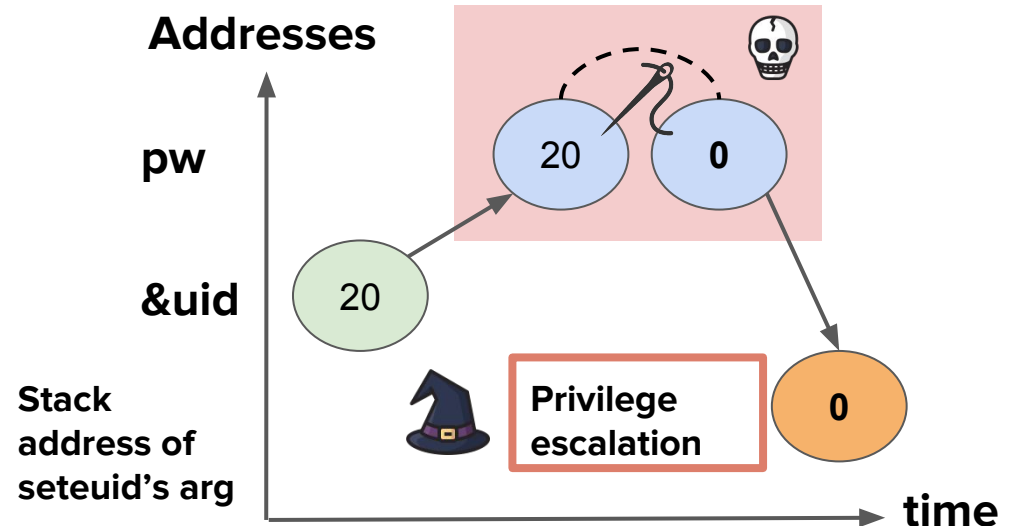
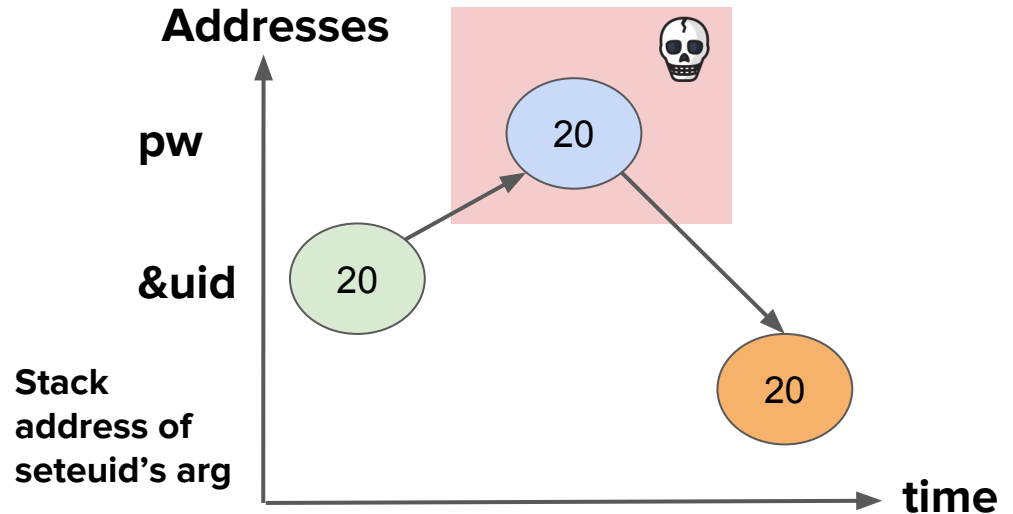
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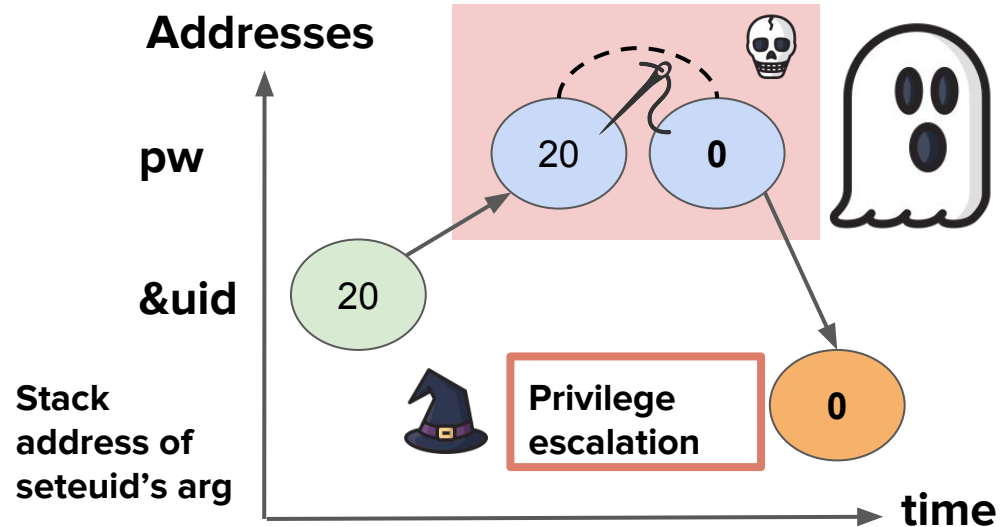
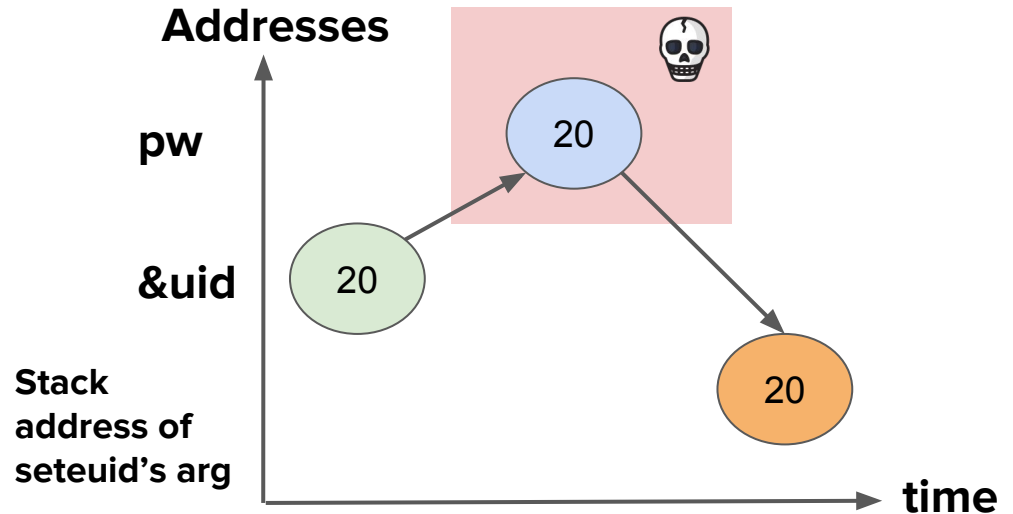
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Data-Oriented Programming (DOP)

Hu et al. Oakland'16

Data-Oriented Programming (DOP)

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Code-reuse attacks with Control-Data

ROP, JOP, SROP

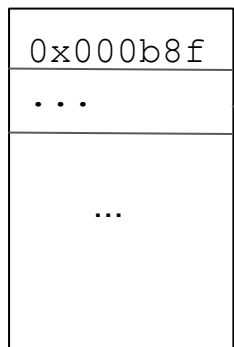
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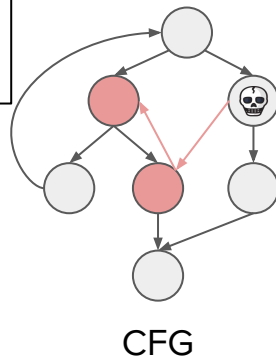
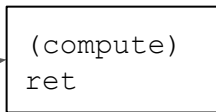
Code-reuse attacks with Control-Data

ROP, JOP, SROP

Addresses of ROP
gadgets



Gadgets



ROP payload

Data-Oriented Programming (DOP)

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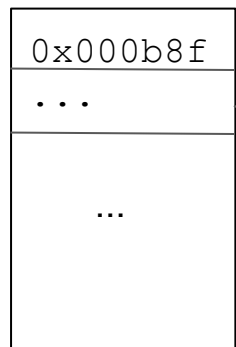
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Requirements:

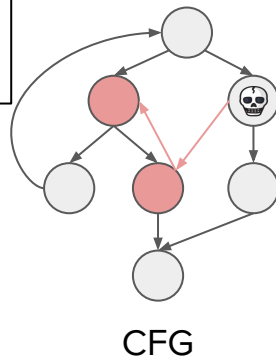
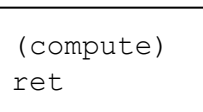
1. Classic gadgets
2. The gadgets must be chained

Addresses of ROP
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ROP payload

Gadgets



CFG

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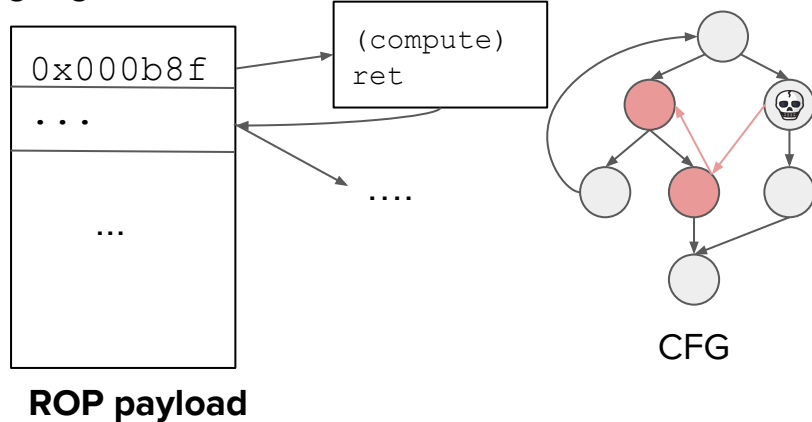
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Code-reuse attacks with Non-Control-Data

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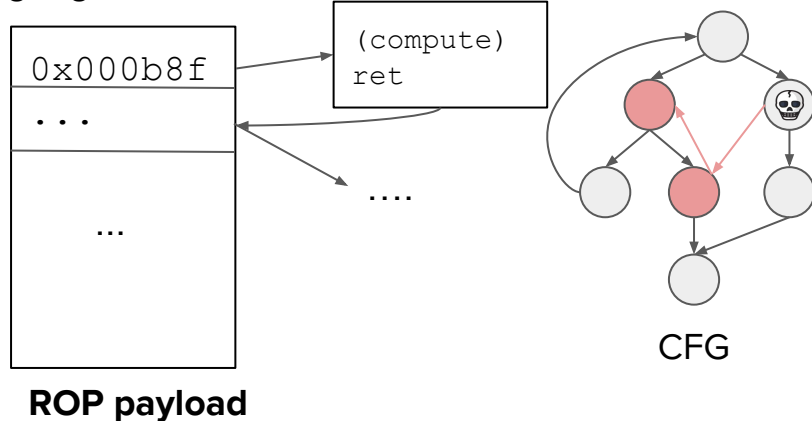
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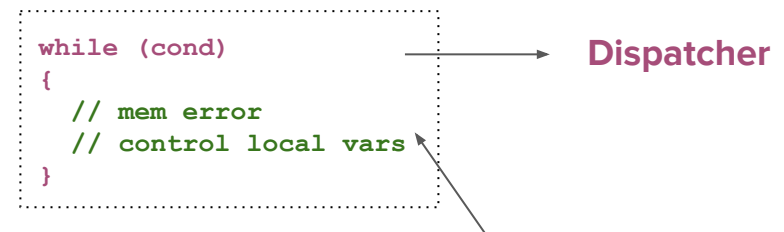
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Addresses of ROP gadgets



Code-reuse attacks with Non-Control-Data

DOP



Operations that change the program's logic
(Data-oriented gadgets)

Data-Oriented Programming (DOP)

Hu et al. Oakland'16

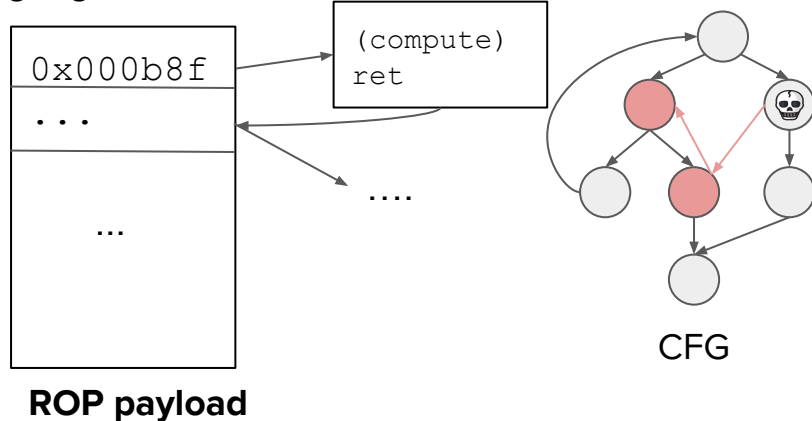
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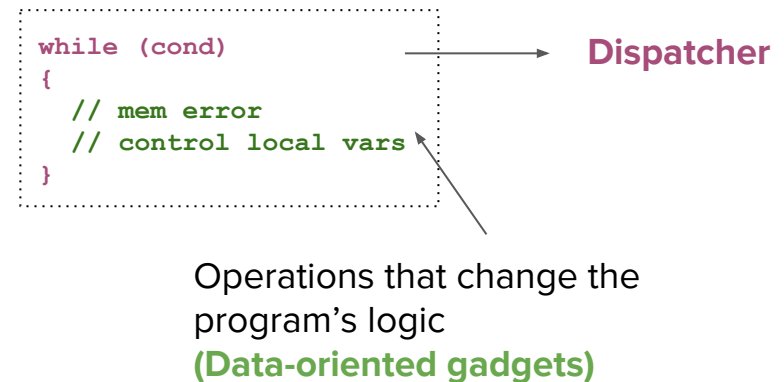


Code-reuse attacks with Non-Control-Data

DOP

Requirements:

1. Data-oriented gadgets
2. Gadget dispatcher



Data-Oriented Programming (DOP)

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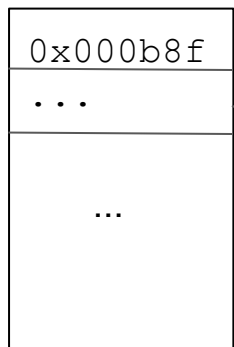
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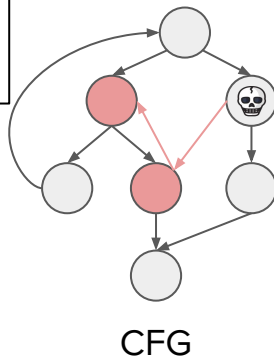
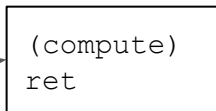
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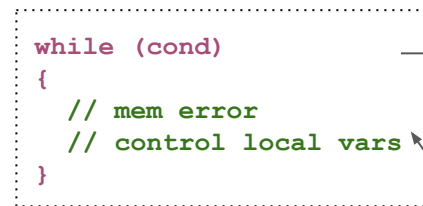
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Code-reuse attacks with Non-Control-Data

DOP

Requirements:

1. Data-oriented gadgets
2. Gadget dispatcher
3. **Must follow the legitimate execution path**



Dispatcher

Operations that change the program's logic
(Data-oriented gadgets)

Data-Flow Integrity (DFI)

Castro et al. OSDI'06

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1 - Offline: DFG computation

Data-Flow Integrity (DFI)

Castro et al. OSDI'06



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a = 0;  
if (...)  
{  
    x = 0;  
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y = x * 2;
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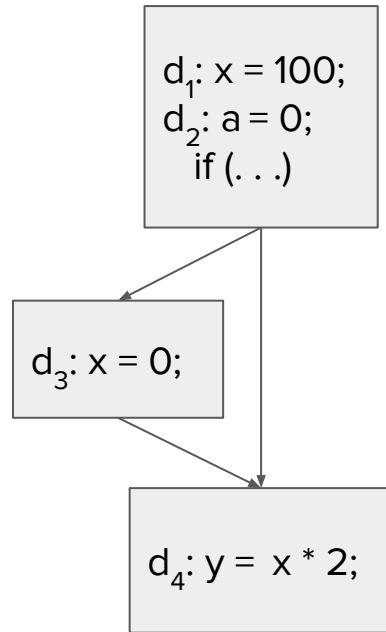
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Data Flow Graph (DFG)

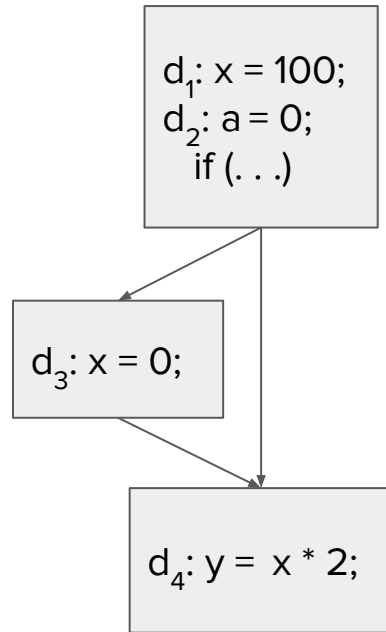
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Data Flow Graph
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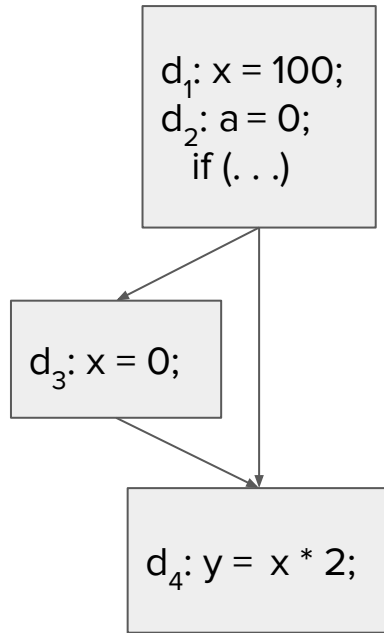
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Data Flow Graph (DFG)

```
x → { d1, d3 }  
a → { d2 }  
y → { d4 }
```

Enforced DFG

1 - Offline: DFG computation

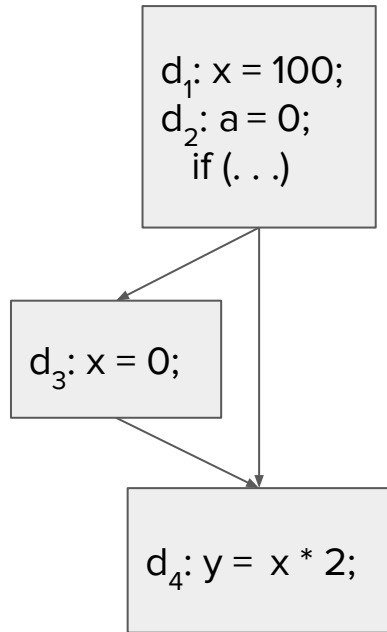
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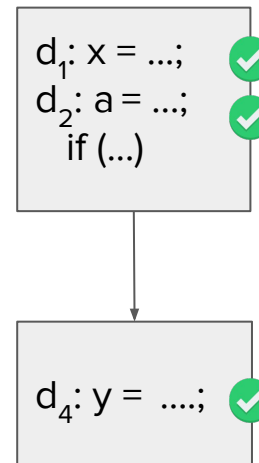


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Enforced DFG

Execution 1



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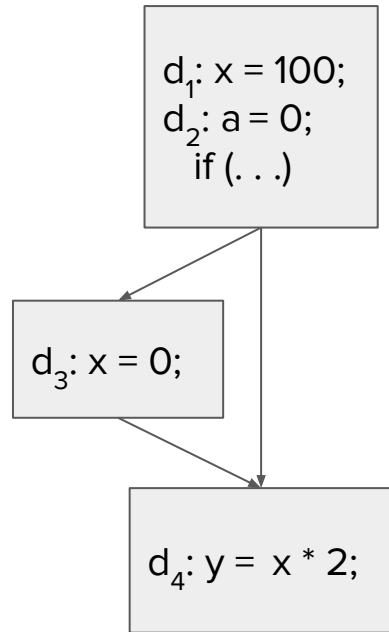
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Data Flow Graph (DFG)

```
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```

Enforced DFG



✓ Execution 1

```
d1: x = ...; ✓  
d2: a = ...; ✓  
if (...)
```

```
d4: y = ...; ✓
```

1 - Offline: DFG computation

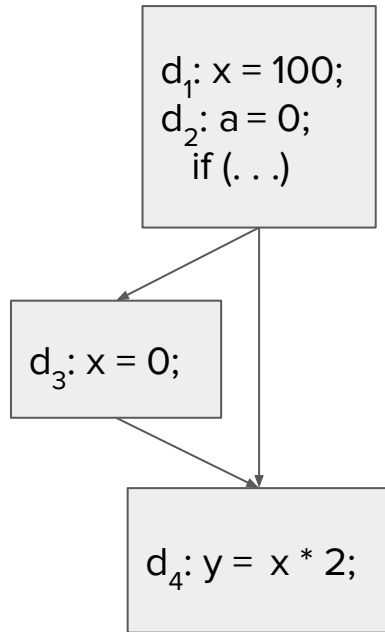
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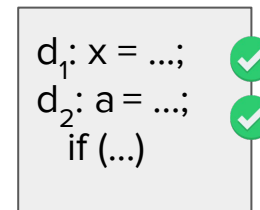


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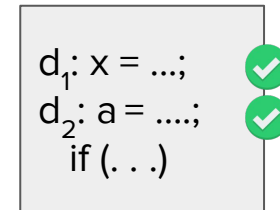
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Enforced DFG

✓ Execution 1



Execution 2



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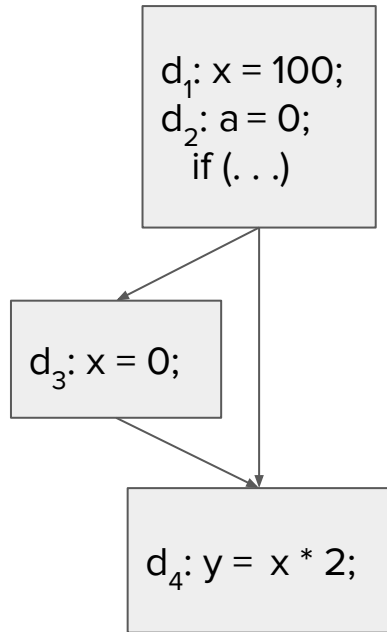
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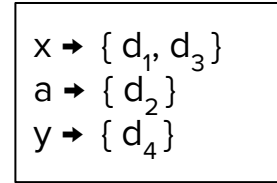
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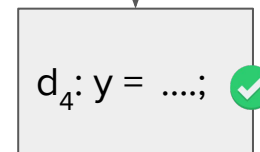
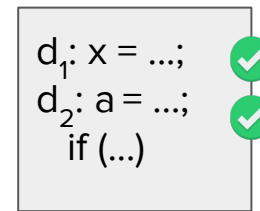


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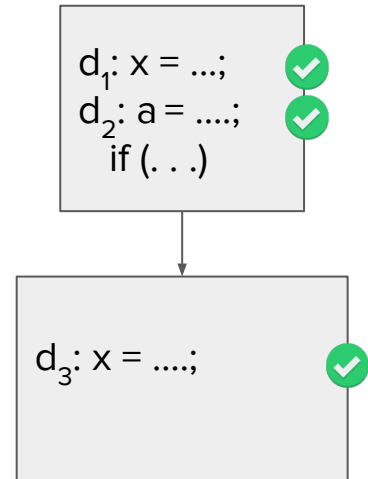


Enforced DFG

✓ Execution 1



Execution 2



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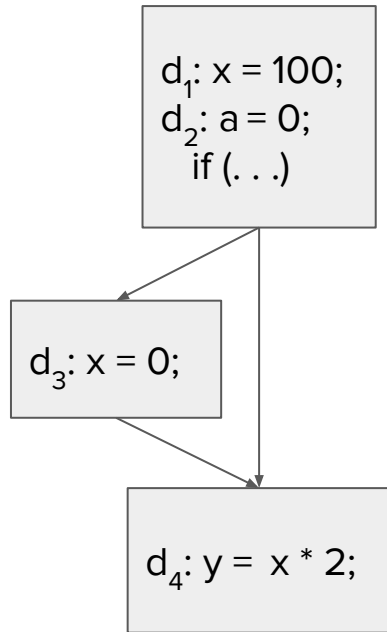
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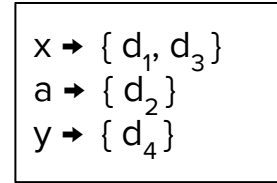
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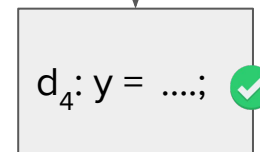
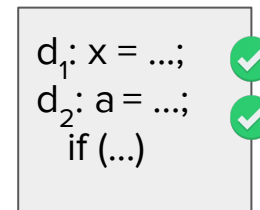


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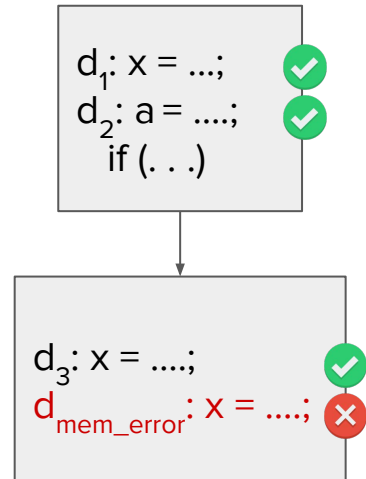


Enforced DFG

Execution 1



Execution 2



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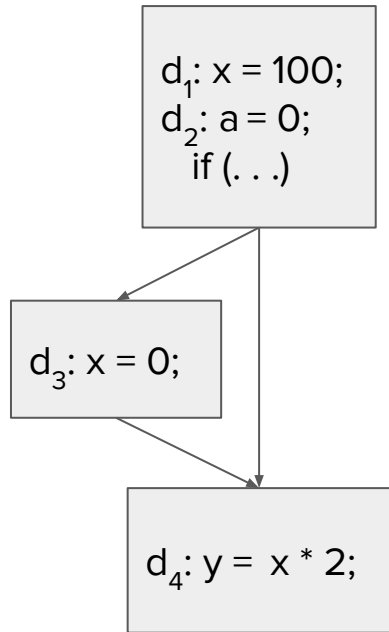
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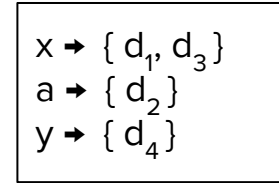
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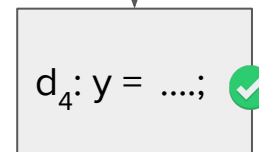
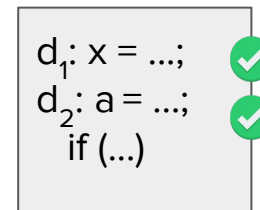
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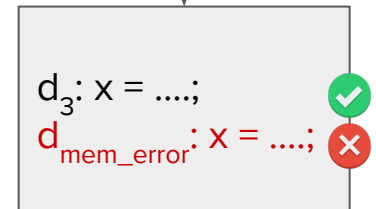
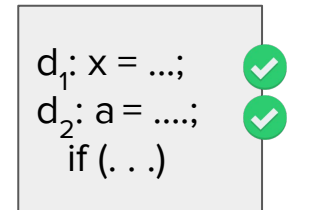
Enforced DFG



✓ Execution 1



✗ Execution 2



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Song et al. NDSS'16

Kernel Data-Flow Integrity

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Protect the kernel against memory-corruption-based privilege escalation attacks

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Protect the access control mechanisms

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Access control checks



Integrity of the code & data of the access controls

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Control-Data + Non-Control-Data



Which are essential to enforce the security invariants?

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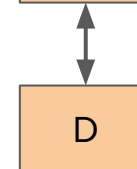
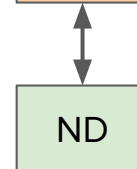
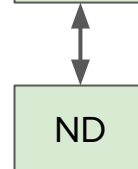
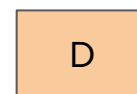
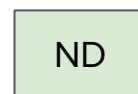
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D: distinguishing, **ND**:non-distinguishing

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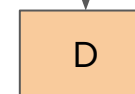
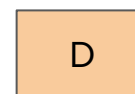
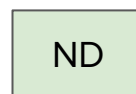
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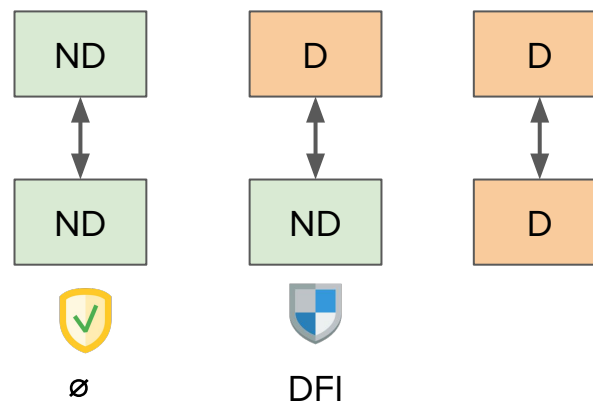
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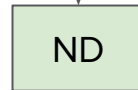
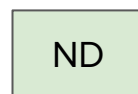
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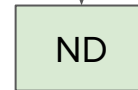
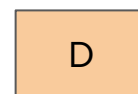


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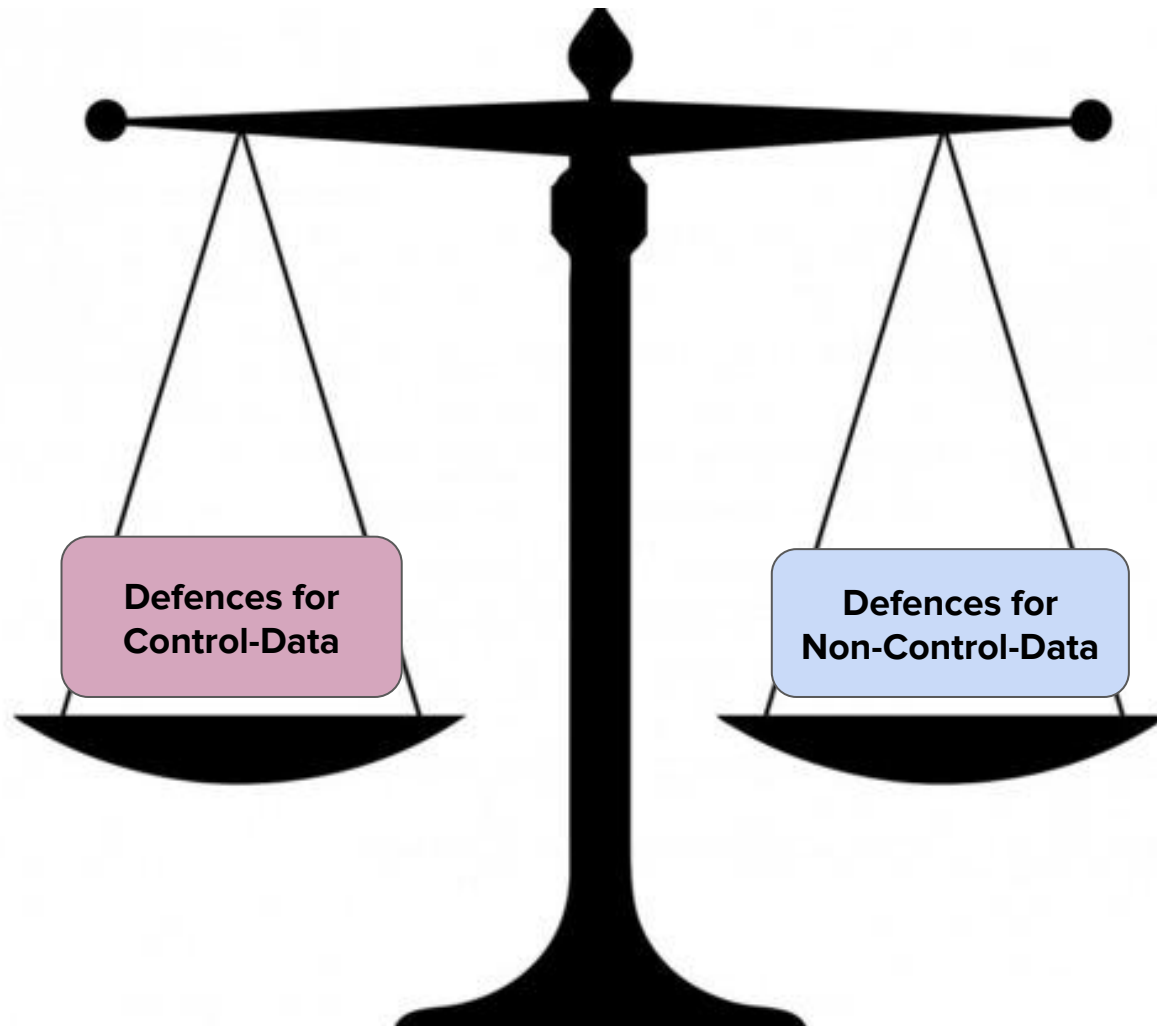
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Closing remarks

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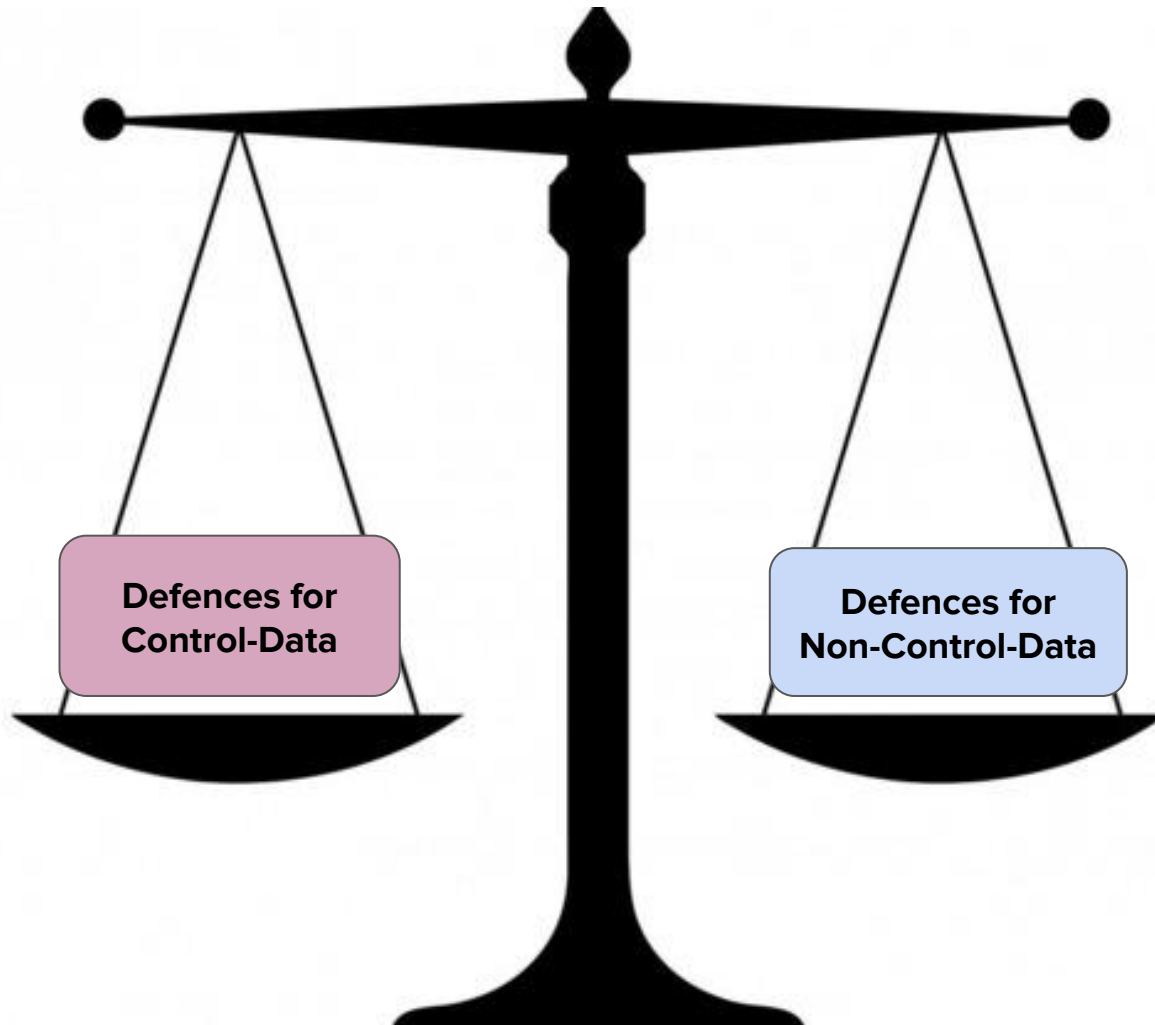


Closing remarks



Why are they
constantly making
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Closing remarks

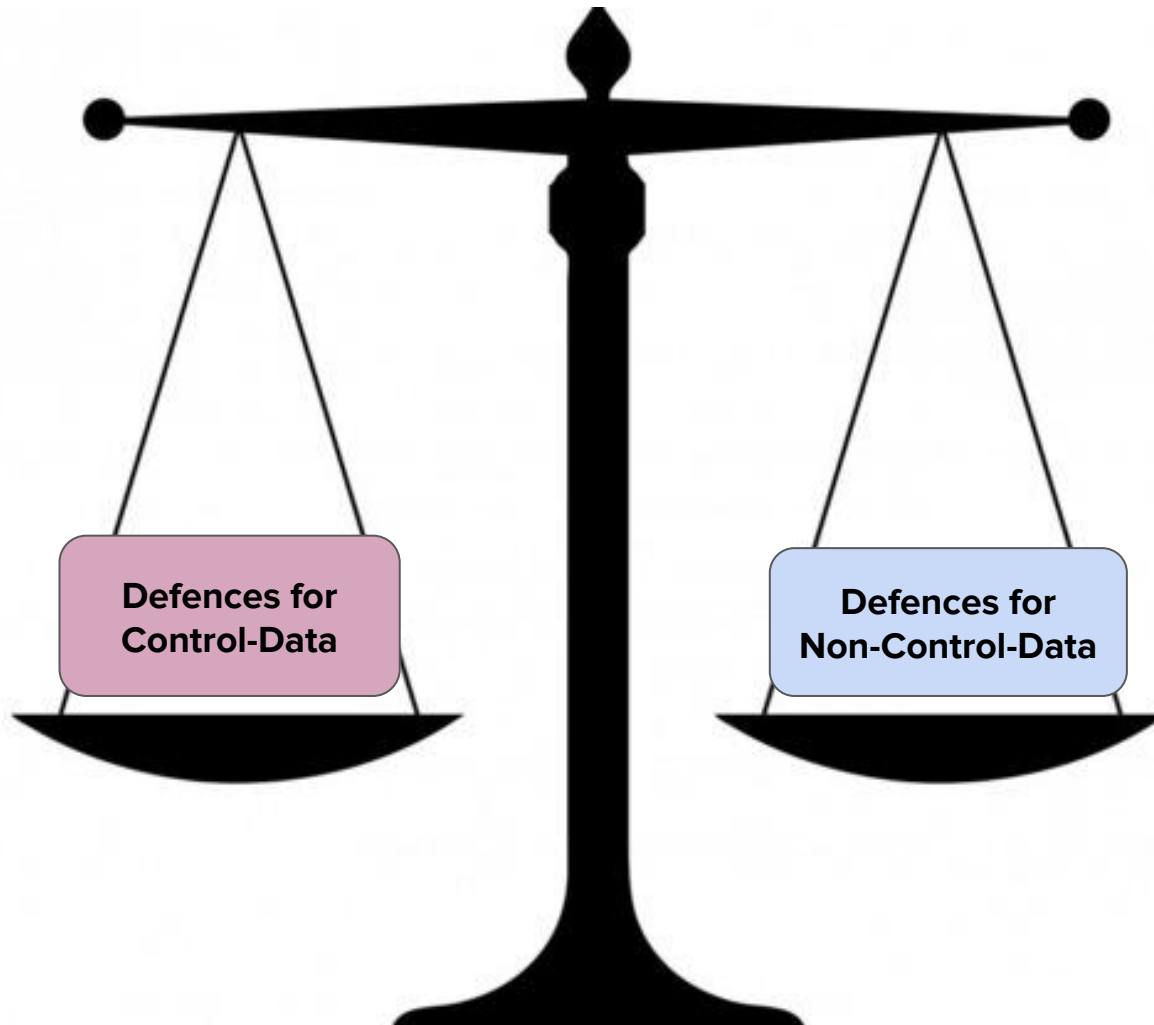


Closing remarks

Why are they constantly making my life harder?



Haters gonna hate



**Defences for
Control-Data**

**Defences for
Non-Control-Data**

Data is Flowing in the Wind: A Review of Data-Flow Integrity Methods to Overcome Non-Control-Data Attacks

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